

Solution by Team Embryo (Third Place) for the SIGMOD Programming Contest 2025 Hangrui Zhou, Yiming Qiao, Shaoxuan Tang



Task: Implement an In-memory Join Pipeline Executor and Operators

- Build a high-performance join pipeline executor that processes in-memory base tables using a given join plan.
- The solution must demonstrate strong end-to-end performance across heterogeneous hardware platforms, including Intel, AMD, ARM, and Power.
- Implementations can leverage multi-threading, SIMD, join reordering, Bloom filters, and other hardware-aware optimizations – but must use the same codebase for all platforms.
- The final output is a fully materialized in-memory result table, and only the total runtime is measured.

Solution: Compact, Clean, and Competitive – Just 2300 Lines of Code

- Vectorized & Push-Based Execution: Efficient tuple processing with high CPU utilization and cache locality.
- Custom Hash Table: Inspired by DuckDB, using salt-based hashing and lightweight Bloom filters for fast key matching.
- Min-Max Filter Pushdown: Early filtering for range conditions to skip irrelevant data efficiently.
- Data Chunk Compaction: Reduce data movement and processing cost by compacting sparse vectors [Qiao, Zhang].
- Lock-Free Multi-threading: Fast and contention-free parallelism across pipeline stages.
- Lean Memory Management: Minimal allocation overhead with simple, allocator-aware design.

Vectorized & Push-Based Execution







- Processes data in batches to reduce per-tuple overhead.
- Minimizes function calls by operating on entire vectors. Wellsuited to modern hardware, leveraging wide SIMD units.
- Ref: MonetDB/X100: Hyper-Pipelining Query Execution. CIDR'05.

Data Chunk Compaction



 Combines linear probing and key-based chaining to handle collisions effectively. On collision, inserts into the next free slot and links it to the same-key chain.

Ref: Adaptive Factorization Using Linear-Chained Hash Tables. CIDR'25.

Implicit SIMD Optimization



- Filters and joins often produce Small chunks.
- It leads to:
 - X High per-chunk overhead; X Poor SIMD/cache utilization.
- Solution: Logical Compaction.
- Ref: Data Chunk Compaction in Vectorized Execution. SIGMOD'25.
- We do not use explicit SIMD intrinsics (e.g., AVX, NEON).
- Instead, our vectorized and tight loop structure allows modern compilers to auto-vectorize.
- Delivers performance gains without sacrificing portability.
- Tight loop also aligns with Group Prefetch.
- Ref: Improving Hash Join Performance through Prefetching. ICDE'04.

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