

What Works and What Does Not: A Winning Strategy for Join Query Execution

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Goal

Execute a join query on columnar data as fast as possible.

Challenge

- α-acyclic equi-joins, PostgreSQL query plans
- Every query processes fresh input data, no statistics
- Bespoke input format

R1.X	=	R2.X
R2.Y	Ξ	R4.Y
R3.Z	=	R5.Z
R4.K	=	R5.K

Query predicates & joins



integer & text

Variable length text data

► IMDB dataset small for modern machines, overheads matter

Approach Keep it simple, keep it fast

Planning

- Statistics based on index-based join sampling [Leis et al. 2017]
- ► DP based join ordering [Moerkotte & Neumann 2006]
- ► Incremental query planing ^[Neumann & Galindo-Legaria 2013]

Table scan

- Fast scans and filtering using bitmaps and vectorization
- ► Bloom filters ^[Birler et al. 2024, Schmidt et al. 2021]

Join pipeline

- ► (Pre-)Compiled join pipelines [Neumann 2011]
- ► Chaining hash table with partitioned loads ^[Birler et al. 2024]
- Eager aggregation of duplicates [Birler et al. 2024]





(2) Optimize Join Plan pick the next cheapest pipeline



for tuple in table:

- *if* <u>not</u> <u>tuple.key</u> *in* probeHt:
 - continue
- for partner in probeHt[tuple.key]:
 - targetHt.insert(tuple + partner)

Infrastructure

- Bump memory allocator
- Efficient scheduling of small and large tasks
- Continuous profiling with Perfetto
- Random test query generation for robustness



Source Code: github.com/umbra-db/contest-sigmod2025

(3) Vectorized Table Scan

filter early, read and process less

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(5) Build Join Filters

prepare for future scans



(6) Repeat update statistics & continue with (1)

Evaluation

Execution Time Improvement with Optimizations

Competitor performance shown as reference lines Each optimization step shows speedup vs. previous



3 repeated executions of JOB queries on prefiltered base tables measured on AMD EPYC 9454P (Linux 6.11.0-26).

Findings

(4) Run Join Pipeline probe joins and build next hash table

- ► The general-purpose database system Umbra is roughly as fast as the second-best solution, even when querycompilation and data-decompression times are included. ► A simple yet efficient **hash table** combined with **compiled**
- join pipelines already provides a strong baseline.
- ► The execution is memory bound. **Early filtering** helps as reading less data makes us faster. SIMD does not make a difference as we are not compute bound.
- PostgreSQL query plans are okay but not great. Adaptive query optimization improves runtime by 2×.